

Regular Pumping Lemma

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1 Regular Pumping Lemma

Regular pumping lemma is a powerful tool that helps you show that certain languages are not regular. The statement of the regular pumping lemma is as follows:

If L is a regular language, then there exists an integer m such that given a word w in L , with $|w| \geq m$, w can be decomposed as $w = xyz$, with the following properties:

1. $|xy| \leq m$
2. $|y| > 0$, and
3. for each $i \geq 0$, $xy^iz \in L$.

To use the pumping lemma to show that a given language L is not regular, we proceed by assuming that L is regular so that L satisfies the pumping lemma. Then we select an appropriate w in L and show that w cannot be decomposed as prescribed by the pumping lemma. This contradiction shows that our assumption must be wrong and hence L is not regular.

We will use the language $L = \{w \in \{a, b\}^* : n_a(w) < n_b(w)\}$.

Thus L is the language of those words in $\{a, b\}^*$ that have fewer as than bs .

We first give a direct argument to show that L is not regular. Then we use JFLAP to practice showing the same result.

2 Direct Proof

Assume that L is regular. Then L satisfies the pumping lemma and hence there is an integer m such that the three properties above are satisfied. Select $w = a^m b^{m+1}$. Then $w \in L$ and $|w| \geq m$. By pumping lemma, $a^m b^{m+1} = w = xyz$. By property 1. of the pumping lemma, y must be a string of as . Suppose $y = a^k$. By property 2. of the pumping lemma, $k > 0$. Then the string $xy^2z = xyyz = a^{m+k} b^{m+1}$ is not in L . This is a contradiction to property 3. Hence L must be nonregular.

3 Proof with JFLAP

JFLAP provides a proof by letting you play a game against the computer. The idea of the game is to win by providing a word w for which pumping lemma fails. Run JFLAP, select

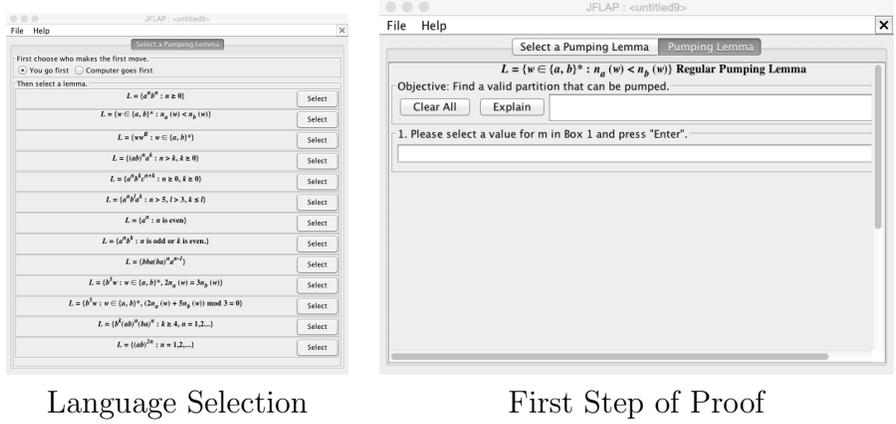


Figure 1: Step 1

Regular Pumping Lemma, select “You go first” and then select the language L as above from the list provided. See Figure 1. Click on “Explain” button to see a description of why L is nonregular. Since you go first, JFLAP asks you to enter a number m . Enter 4. See Figure 2.

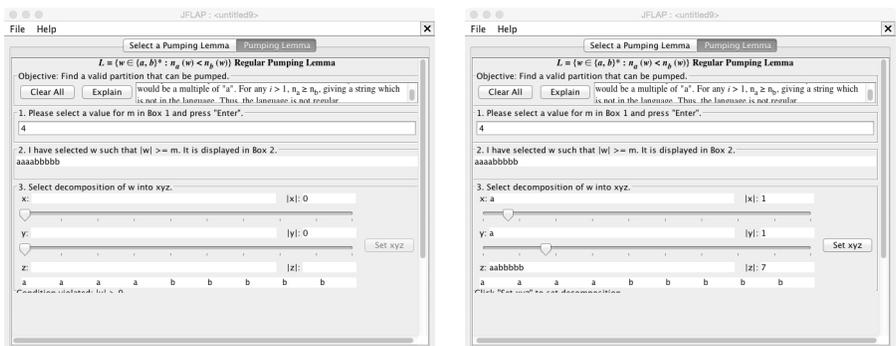


Figure 2: Step 2

At this point JFLAP gives you the word $w = a^4b^5$ and asks you to provide a decomposition of w . Use the sliders to select x and y . Because of properties 1. and 2. of the pumping lemma, y must be one of $a, aa, aaa, aaaa$. Select $y = aa$ and then select $x = a$. z is then selected for you. If x and y are selected appropriately, then the button “Select xyz ” can be clicked. Click it. JFLAP then presents you with $i = 2$ and shows that the string $xyyz$ is not in L . You can animate that proof. See Figure 3. Repeat this entire process for some other values of m .

We now do this again but this time we’ll let the “Computer go first.” Suppose it selects 15 as in Figure 4. You are asked to select a w . Make sure that w is in L and $|w| \geq 15$. What is an appropriate w ? See Figure 4. With the right choice of w and then an i , you win. Try choosing different w s. Repeat the entire process two or three times.

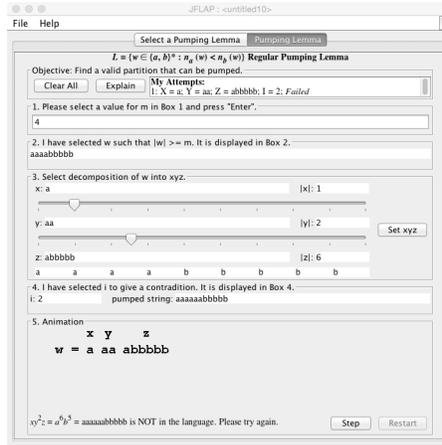


Figure 3: Step 3

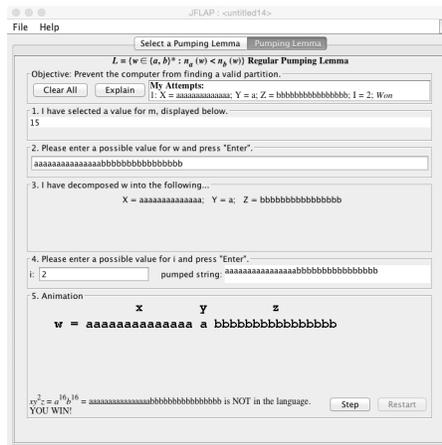


Figure 4: Computer goes first

4 References

1. Introduction to the Theory of Computation (Third Edition), Michael Sipser. Cengage Learning. 2013.
2. JFLAP - An Interactive Formal Languages and Automata Package, Susan H. Rodger and Thomas W Finley. Jones and Bartlett Publishers. 2006